

# Some properties of epistemic models expressible in group announcement logic are not expressible in coalition announcement logic



## Groups versus Coalitions: On the Relative Expressivity of GAL and CAL

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### 1. Preliminaries

Public announcement logic (PAL) [3] is used to reason about the effects of truthful public announcements on agents' knowledge. Group announcement logic (GAL) [1] and coalition announcement logic (CAL) [2] extend PAL with quantification over announcements. In our work, we show that there are some properties of models that are expressible in GAL and are not expressible in CAL.

**Definition 1 (Epistemic Model)** An epistemic model is a triple  $M = (S, \sim, V)$ , where

- $S$  is a non-empty set of states,
- $\sim_i \subseteq S^2$  is an equivalence relation to each agent  $i$ ,
- $V : P \rightarrow 2^S$  is the valuation function.

A pair  $M_s$  with  $s \in S$  is called a pointed model.

An announcement of  $\alpha$  in a pointed model  $M_s$  results in an updated pointed model  $M_s^{\|\alpha\|_M}$  containing only  $\alpha$ -states:

- $S^{\|\alpha\|_M} = \|\alpha\|_M$ ,
- $\sim_i^{\|\alpha\|_M} = \sim_i \cap (\|\alpha\|_M \times \|\alpha\|_M)$ ,
- $V^{\|\alpha\|_M}(p) = V(p) \cap \|\alpha\|_M$ .

**Definition 2 (Semantics)** Let  $M_s$  be a pointed epistemic model. The semantics for boolean cases is as usual.

$$M_s \models K_i \phi \text{ iff } \forall t \in S : s \sim_i t \text{ implies } M_t \models \phi$$

$$M_s \models [\alpha] \phi \text{ iff } M_s \not\models \alpha \text{ or } M_s^{\|\alpha\|_M} \models \phi$$

$$M_s \models \langle \alpha \rangle \phi \text{ iff } M_s \models \phi \text{ and } M_s^{\|\alpha\|_M} \models \phi$$

$[\alpha] \phi, \langle \alpha \rangle \phi$ : 'after public announcement of  $\alpha$ ,  $\phi$  holds in the resulting model'.

### 2. Announcements by groups and coalitions

We are interested in the following restrictions on announcements:

- Announcements are made by agents
- Agents can only announce what they know
- Groups of agents can announce conjunctions of formulas, where each conjunct is a formula known by an agent in the group
- Coalitional ability wrt announcements: can a group make an announcement such that whatever agents outside the group announce, the goal formula holds?

**Group Announcement Logic (GAL)** = PAL +  $\{[G]\phi, \langle G \rangle \phi\}$

$\langle G \rangle \phi$ : 'there is an announcement by agents from  $G$  such that  $\phi$  holds in the resulting model'

$[G]\phi$ : 'whatever agents from  $G$  announce,  $\phi$  holds in the resulting model'

**Coalition Announcement Logic (CAL)** = PAL +  $\{\langle G \rangle \phi, \langle G \rangle \phi\}$

$\langle G \rangle \phi$ : 'there is an announcement by agents from  $G$  such that whatever agents  $A \setminus G$  outside of the coalition announce,  $\phi$  holds'

$\langle G \rangle \phi$ : 'whatever agents from  $G$  announce, there is an announcement by the agents from the outside of the coalition, such that  $\phi$  holds'

Let  $\alpha_G$  be a shorthand for a formula of the type  $K_i \phi_i \wedge \dots \wedge K_j \phi_j$ , where  $i, \dots, j \in G$ , and  $\phi_i, \dots, \phi_j$  are formulas of epistemic logic.

$$M_s \models [G]\phi \text{ iff } \forall \alpha_G : M_s \models [\alpha_G] \phi$$

$$M_s \models \langle G \rangle \phi \text{ iff } \forall \alpha_G \exists \beta_{A \setminus G} : M_s \models \alpha_G \rightarrow \langle \alpha_G \wedge \beta_{A \setminus G} \rangle \phi$$

$$M_s \models \langle G \rangle \phi \text{ iff } \exists \alpha_G : M_s \models \langle \alpha_G \rangle \phi$$

$$M_s \models \langle G \rangle \phi \text{ iff } \exists \alpha_G \forall \beta_{A \setminus G} : M_s \models \alpha_G \wedge [\alpha_G \wedge \beta_{A \setminus G}] \phi$$

### 3. Groups versus coalitions

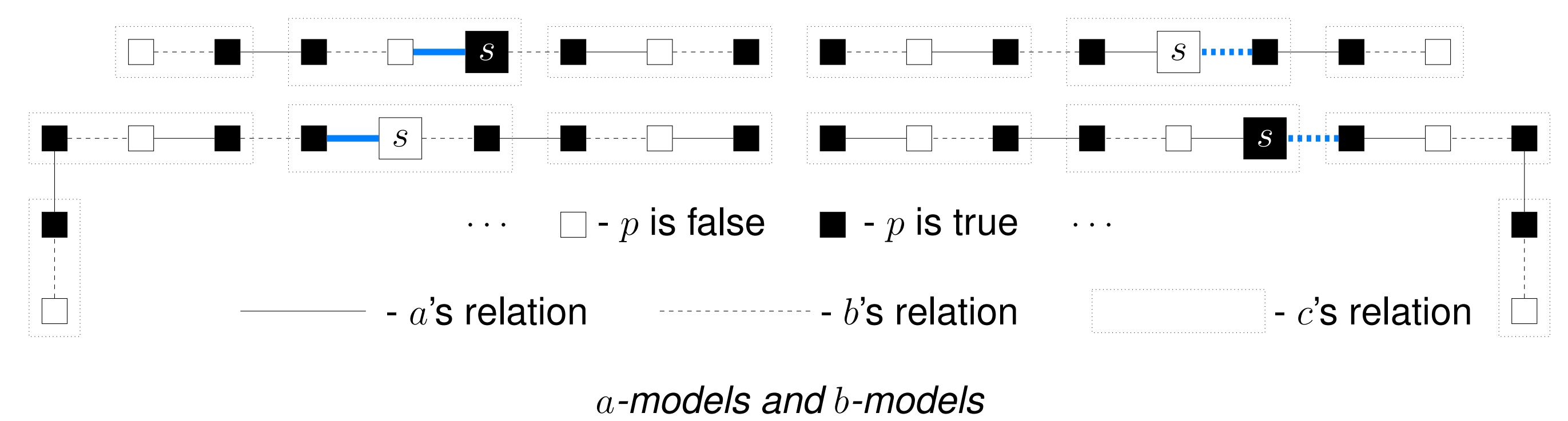
**Definition 3 (Expressive power)** Let  $\mathcal{L}_1$  and  $\mathcal{L}_2$  be two languages.  $\mathcal{L}_1$  is at least as expressive as  $\mathcal{L}_2$  ( $\mathcal{L}_2 \leq \mathcal{L}_1$ ) iff for all  $\phi \in \mathcal{L}_2$  there is a  $\psi \in \mathcal{L}_1$  such that  $\phi$  and  $\psi$  are satisfied in the same pointed models.

**Theorem 1** GAL  $\not\leq$  CAL

One ingredient in the proof are *formula games*:

- A game is between the  $\forall$ -player and the  $\exists$ -player
- If either player cannot move, she loses
- Coalition announcements  $\langle G \rangle \phi$  and  $\langle G \rangle \phi$  are split into announcements by the coalition and announcements by the anti-coalition using half-coalition announcements  $\langle A \setminus G, \psi_G \rangle$  and  $\langle A \setminus G, \psi_G \rangle$

To show that a property  $\phi$  is expressible in GAL but not CAL, we have to work with infinite classes of models (because both CAL and GAL include PAL, and in a given model a fixed PAL formula can be used for any announcement).



- Two infinite classes of models  $\mathcal{M}_A$  and  $\mathcal{M}_B$
- We define the property  $\Phi$  of being an  $a$ -model expressible in GAL, i.e. for all  $M_s \in \mathcal{M}_A$ :  $M_s \models \Phi$  and for all  $N_t \in \mathcal{M}_B$ :  $N_t \not\models \Phi$
- Using formula games, we argue that for all  $\Psi \in \mathcal{L}_{CAL}$ , if for all  $M_s \in \mathcal{M}_A$ :  $M_s \models \Psi$ , then there is an  $N_t \in \mathcal{M}_B$ :  $N_t \models \Psi$
- CAL operators require all agents to make a simultaneous announcement. For chain models, the intersection of all agents' relations is the identity, and hence if a coalition can force some configuration of an  $a$ -model, then it can replicate the same configuration in a  $b$ -model. GAL operators, however, do not require all agents to participate in an announcement, and relations of individual agents are not discerning enough to force isomorphic submodels of given  $a$ - and  $b$ -models.

### References

- [1] Thomas Ágotnes, Philippe Balbiani, Hans van Ditmarsch, Pablo Seban. Group announcement logic. *Journal of Applied Logic*, 8(1), 2010, pages 62–81.  
 [2] Thomas Ágotnes and Hans van Ditmarsch. Coalitions and Announcements. In *Proceedings of AAMAS 2008*, pages 673–680.  
 [3] Jan Plaza. Logics of public communications. *Synthese* 158(2), 2007, pages 165–179.

